



Middlebury

CSCI 201: Data Structures

Fall 2024

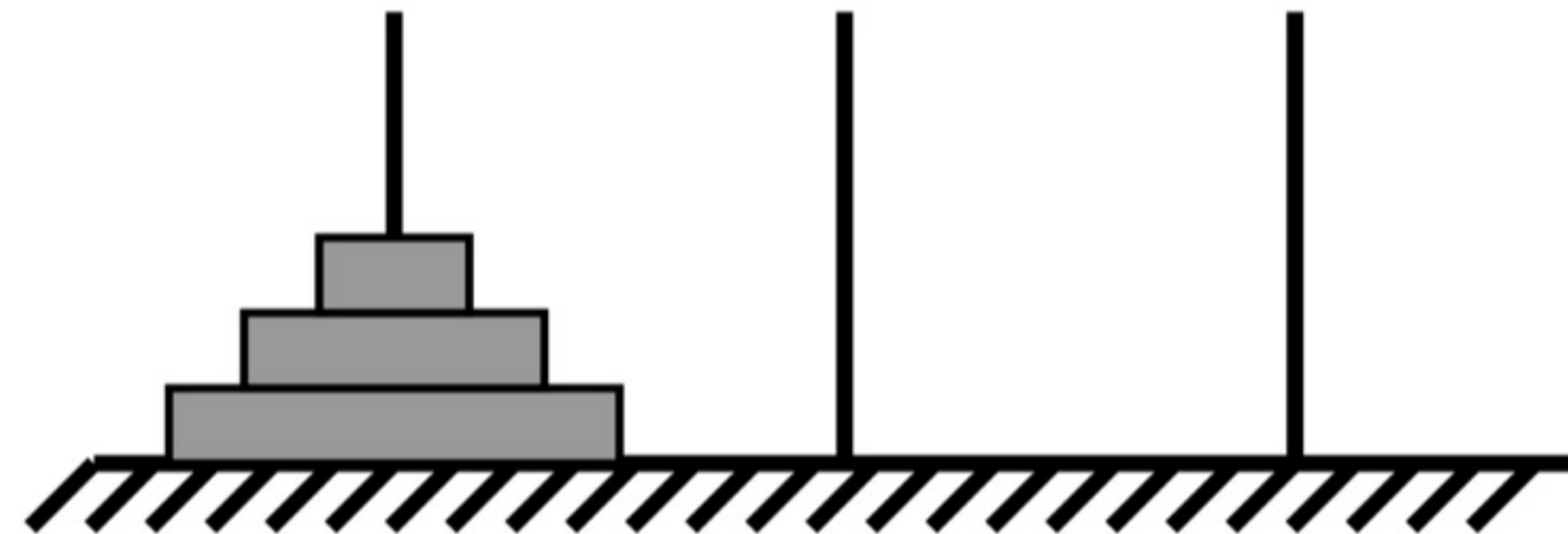
Lecture 5T: Recursion

Goals for today:

- Associate a call to a method with a **frame**.
- Use **recursion** to solve problems by building a solution from solutions to smaller problems.
- Ensure we always have a **base case** in our recursive solutions.
- Use **tail recursion** so that there are no pending operations after the recursive call.
- Practice some more with big-oh notation.

→ + approaching the base case.

$n = 3$ 7
 $n = 4$ 15
 $n = 5$ 31



how many "moves"
to displace stack
of n disks?
($n = 3$, here)

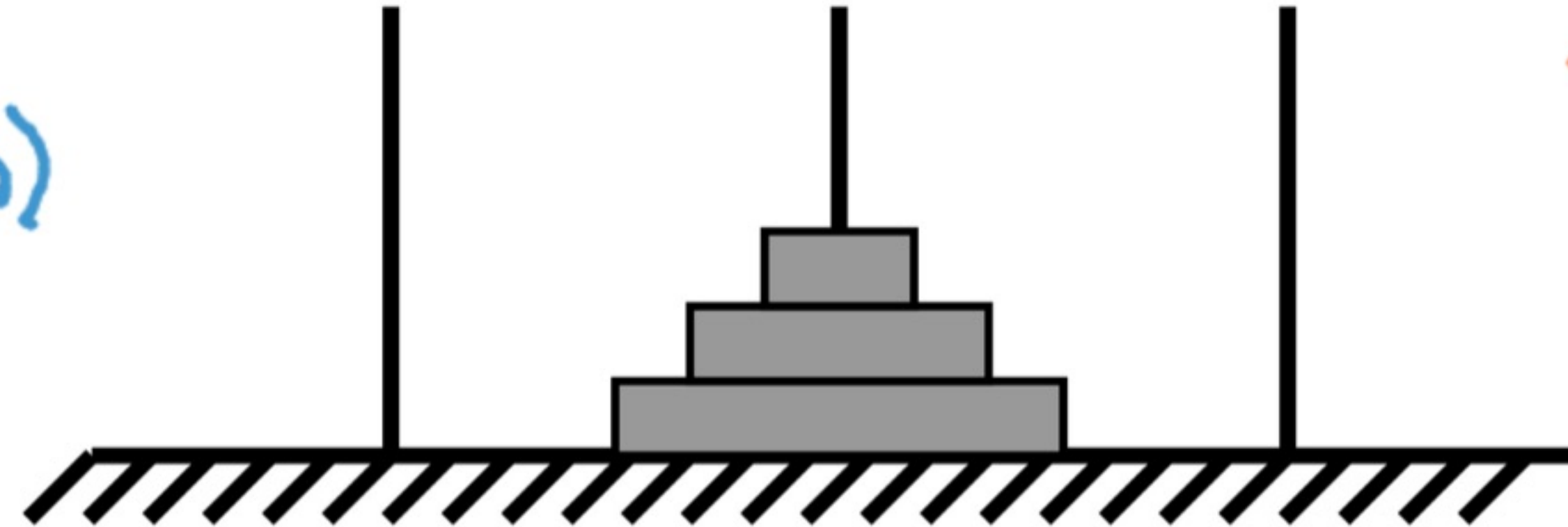
Rules: (also play here! <https://www.mathsisfun.com/games/towerofhanoi.html>)

1. Move one disk at a time.
2. Every move displaces a disk from the top of one stack to the top of another (or empty rod).
3. No larger disk can ever be placed on top of a smaller one.

Counting the number of "moves" in the Tower of Hanoi problem.

think of "moves"
like operations
(last week)

$T(n)$

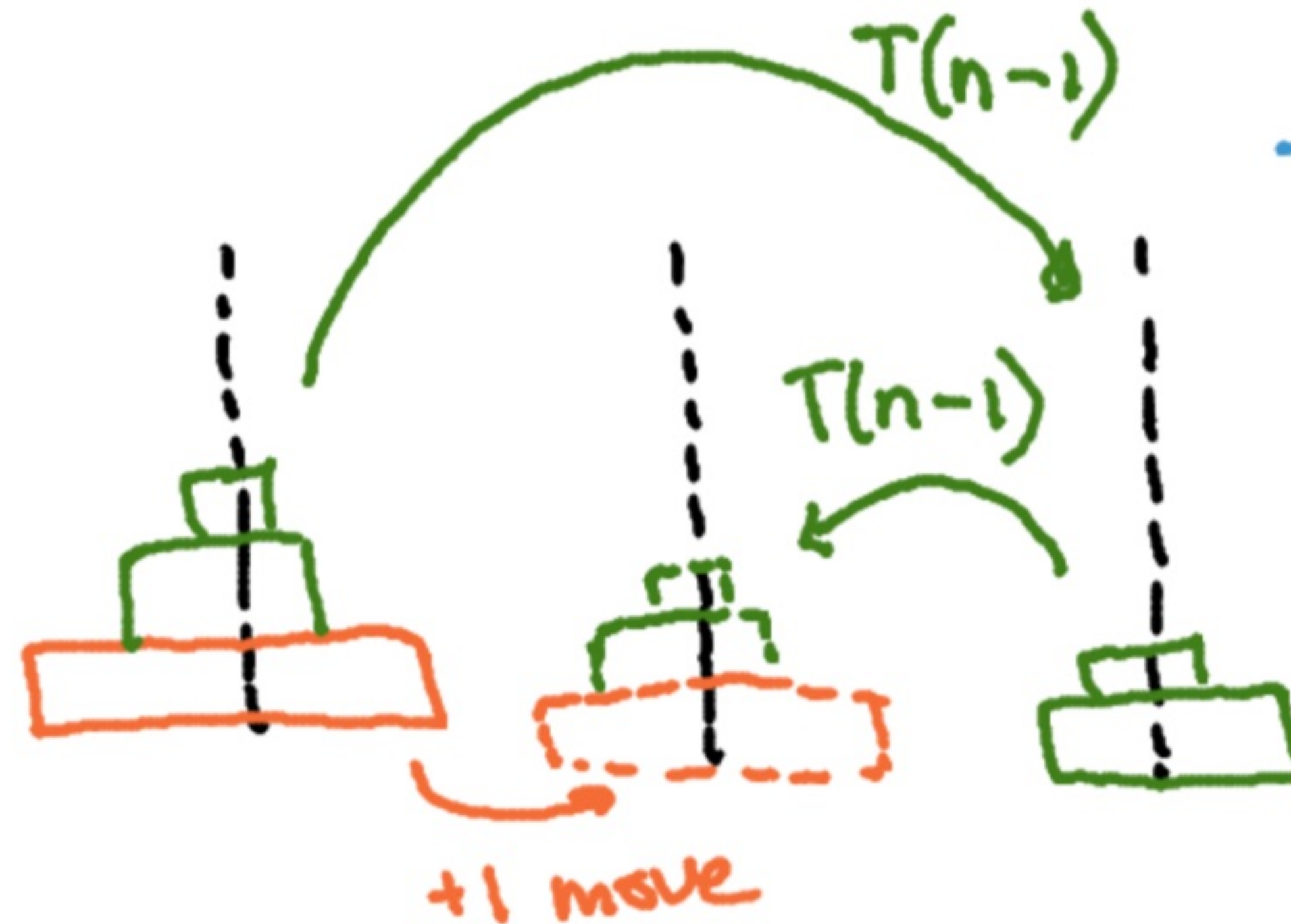


~~||||~~ || 7

recursive

$$T(n) = 2T(n-1) + 1$$

$$T(1) = 1 \text{ (base case)}$$



Two ingredients for a recursive method:

- **Base case:** represents a problem you know how to solve.
- **Recursive case:** builds the solution to a problem from smaller subproblems.

Arithmetic series:

$$1 + 2 + 3 + \dots + (n - 1) + n$$

$$= \frac{n(n+1)}{2}$$

```
1 public static int sum(int n) {  
2     if (n == 1) { // base case  
3         return 1;  
4     }  
5     // recursive case  
6     return n + sum(n - 1);  
7 }
```

approach $n == 1$

Factorial:

$$1 \cdot 2 \cdot 3 \cdot \dots \cdot (n - 1) \cdot n$$

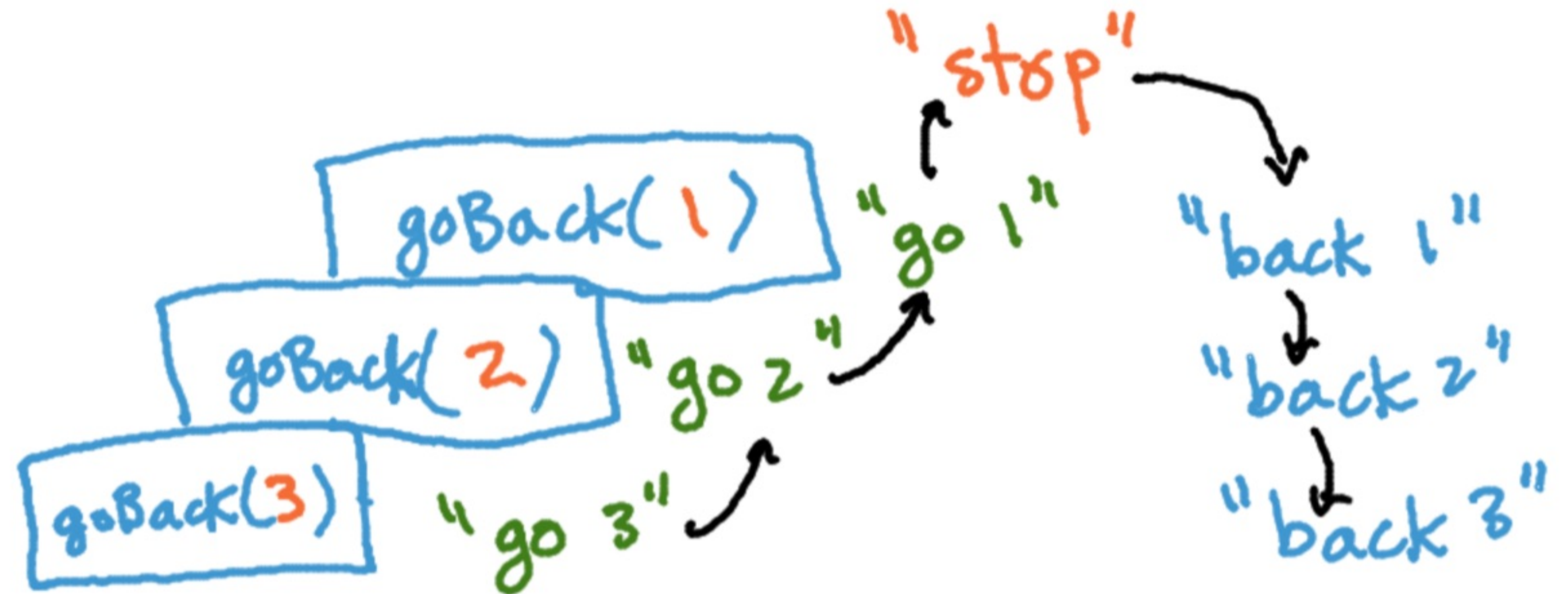
```
1 public static int factorial(int n) {  
2     if (n < 2) { // base case  
3         return 1;  
4     }  
5     // recursive case  
6     return n * factorial(n - 1);  
7 }
```

approach $n < 2$.

The recursive case needs to make progress towards the base case!

What does the **Call Stack** look like?

Open **GoBackExample.java** and add breakpoint on Line 6.



```
1 public static void goBack(int n) {  
2     if (n == 0) {  
3         System.out.println("Stop");  
4     } else {  
5         System.out.println("Go " + n);  
6         goBack(n - 1);  
7         System.out.println("Back " + n);  
8     }  
9 }
```

recursion (n.)

"return, backward movement," 1610s, from Latin *recursionem* (nominative *recursio*) "a running backward, return," noun of action from past-participle stem of *recurrere* "run back" (see **recur**).

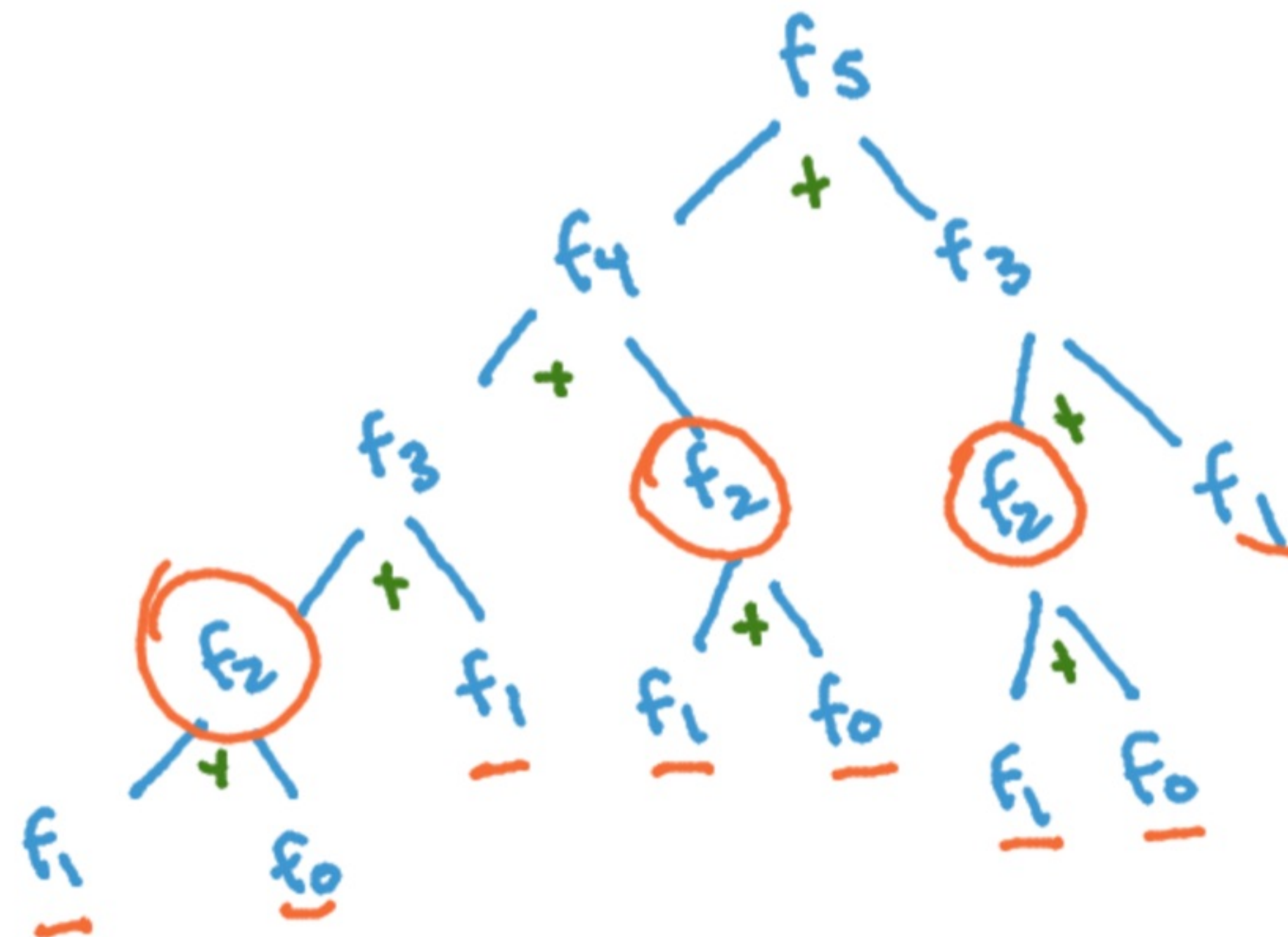
Exercise: write a recursive function to compute Fibonacci numbers.

$$f(n) = f(n - 1) + f(n - 2), \quad \text{with } f(0) = 0, f(1) = 1.$$

```
1 public static int fib(int n) {  
2     if (n < 2) {  
3         return n;  
4     }  
5     return fib(n - 1) + fib(n - 2);  
6 }
```

Try $n = 40, n = 45, n = 48$. What do you notice?

f_2 calculated
3 times



levels $\approx n$

branches = 2

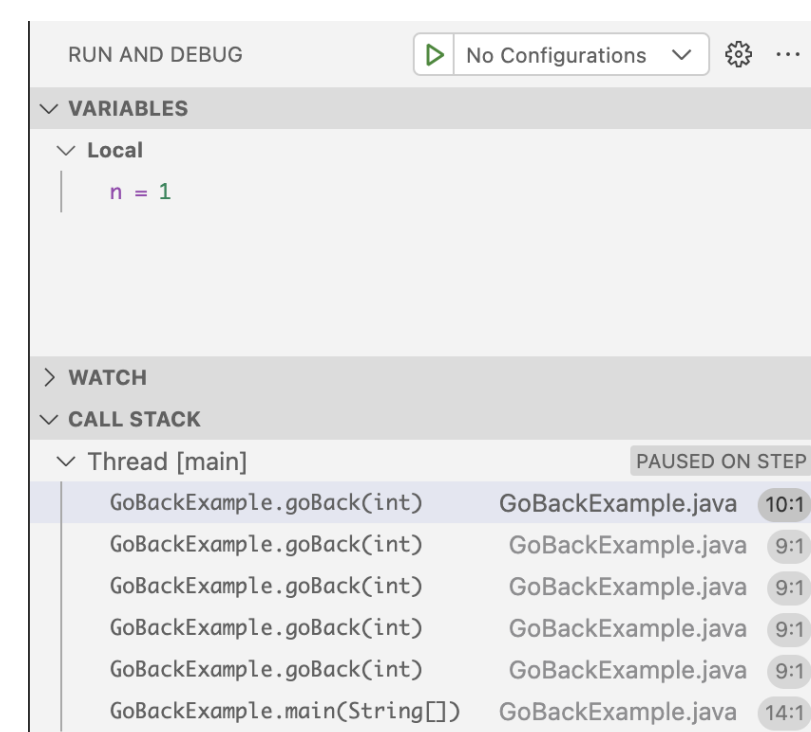
$\approx 2^n$ calls
to fib

$O(2^n)$
exponential
slow 😞

The downsides of recursion.

- Performance might not be great.
- Recursion depth is limited by the stack size.

Go back to **GoBackExample.java** (remove breakpoints) and try $n = 100000$.



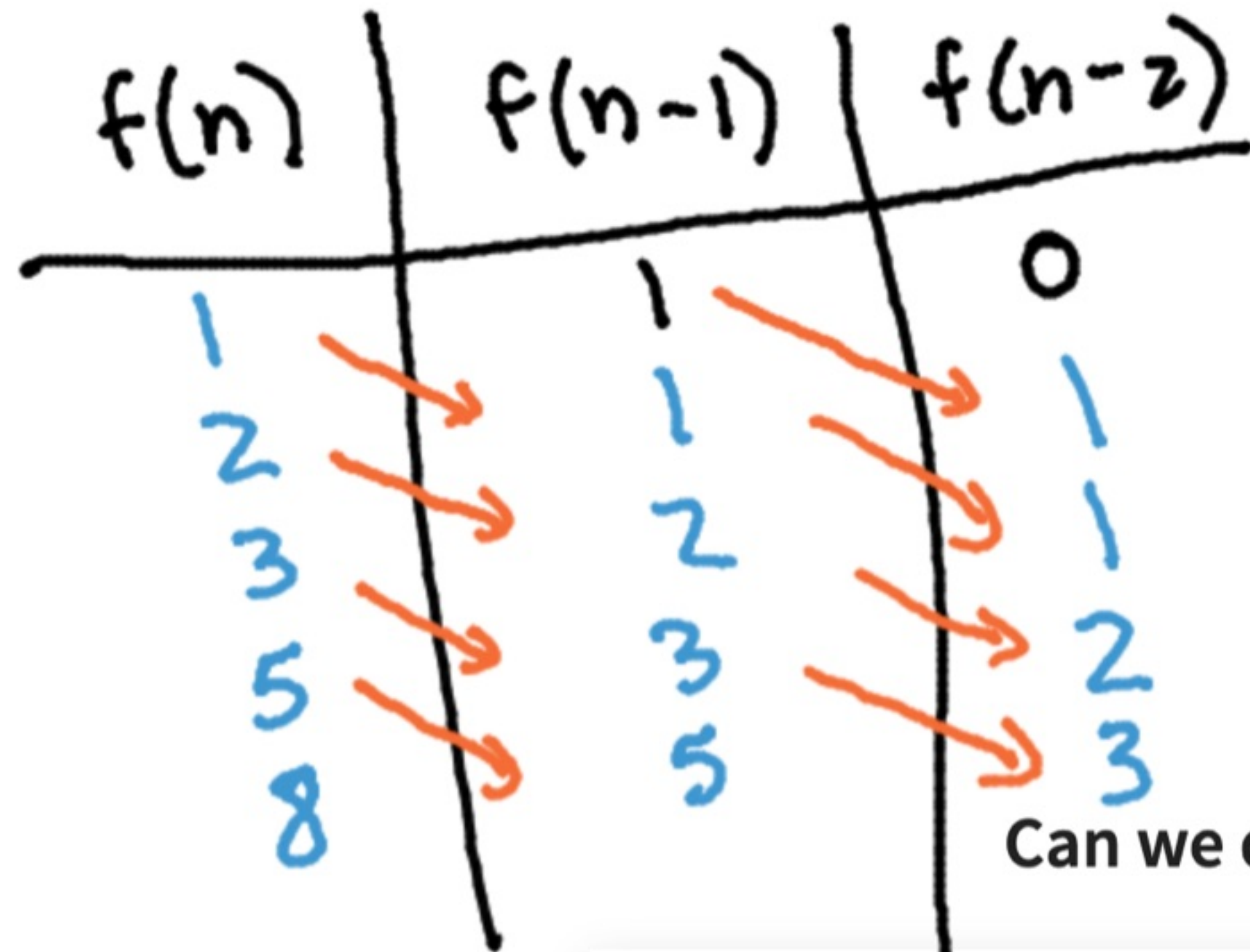
```
1 public static void goBack(int n) {  
2     if (n == 0) {  
3         System.out.println("Stop");  
4     } else {  
5         System.out.println("Go " + n);  
6         goBack(n - 1);  
7         System.out.println("Back " + n);  
8     }  
9 }
```

```
Go  
Go  
Go  
Exception in thread "main" java.lang.StackOverflowError  
at java.base/java.io.FileOutputStream.write(FileOutputStream.java:349)  
at java.base/java.io.BufferedOutputStream.flushBuffer(BufferedOutputStream.java:81)  
at java.base/java.io.BufferedOutputStream.flush(BufferedOutputStream.java:142)  
at java.base/java.io.PrintStream.write(PrintStream.java:570)  
at java.base/sun.nio.cs.StreamEncoder.writeBytes(StreamEncoder.java:234)  
at java.base/sun.nio.cs.StreamEncoder.implFlushBuffer(StreamEncoder.java:313)  
at java.base/sun.nio.cs.StreamEncoder.flushBuffer(StreamEncoder.java:111)  
at java.base/java.io.OutputStreamWriter.flushBuffer(OutputStreamWriter.java:178)  
at java.base/java.io.PrintStream.println(PrintStream.java:723)  
at java.base/java.io.PrintStream.println(PrintStream.java:1028)  
at GoBackExample.goBack(GoBackExample.java:8)  
at GoBackExample.goBack(GoBackExample.java:9)  
at GoBackExample.goBack(GoBackExample.java:9)
```



Can we make our Fibonacci number calculation more efficient?

Let's try to write it as a **while**-loop instead.



```
1 public static int fibWhile(int n) {  
2     int fnMinus2 = 0;  
3     int fnMinus1 = 1;  
4     int i = n;  
5     while (i > 1) {  
6         int fn = fnMinus1 + fnMinus2;  
7         fnMinus2 = fnMinus1;  
8         fnMinus1 = fn;  
9         i--;  
10    }  
11    return fnMinus1;  
12 }
```

Can we do something similar using recursion?

```
1 public static int fib(int n) {  
2     if (n < 2) return n;  
3     return fibHelper(n, 1, 0);  
4 }  
5  
6 private static int fibHelper(int n, int fnMinus1, int fnMinus2) {  
7     if (n == 1) return fnMinus1;  
8     return fibHelper(n - 1, fnMinus1, fnMinus2);  
9 }
```


Can we make our *recursive* Fibonacci number calculation more efficient?

```
1 public static int fibWhile(int n) {
2     int fnMinus2 = 0;
3     int fnMinus1 = 1;
4     int i = n;
5     while (i > 1) {
6         int fn = fnMinus1 + fnMinus2;
7         fnMinus2 = fnMinus1;
8         fnMinus1 = fn;
9         i--;
10    }
11    return fnMinus1;
12 }
```

```
1 public static int fib(int n) {
2     if (n < 2) return n;
3     return fibHelper(n, 1, 0);
4 }
5
6 private static int fibHelper(int n, int fnMinus1, int fnMinus2) {
7     if (n == 1) return fnMinus1;
8     return fibHelper(n - 1, fnMinus1, fnMinus1 + fnMinus2);
9 }
```

fnMinus1 + fnMinus2

≡ cs201-lecture05T



≡ What should the recursive call to fibHelper look like on Line 8? (slido.com #3769190)

33

☐ fibHelper(n - 1, fnMinus1, fnMinus1 + fnMinus2);

☐ fibHelper(n - 1, fnMinus2, fnMinus1 + fnMinus2);

☒ fibHelper(n - 1, fnMinus1 + fnMinus2, fnMinus1);

Send

Voting as Anonymous



Two things going on here:

- Computing any given Fibonacci number *once*.
- **Tail-Recursive:** the recursive call is the **last** thing done in the method.
 - **Strategy to develop:** think about how to write it as a **while**-loop first.

helper

[

```
1 public static int fib(int n) {  
2     if (n < 2) return n;  
3     return fibHelper(n, 1, 0);  
4 }  
5  
6 public static int fibHelper(int n, int fnMinus1, int fnMinus2) {  
7     if (n == 1) return fnMinus1;  
8     return fibHelper(n - 1, fnMinus1 + fnMinus2, fnMinus1);  
9 }
```

How many times is **fibHelper** called?

Same as while loop
 $O(n)$

why tail recursion?

Some compilers will detect it's the last thing called and will optimize out the creation of a ^{new} stack frame.

Exercise: rewrite **factorial** to be tail-recursive.

```
int f = 1;
int i = 1;
while (i < n) {
    f = f * i;
    i++;
}
```

```
1 public static int factorial(int n) {
2     if (n < 2) { // base case
3         return n;
4     }
5     // recursive case
6     return n * factorial(n - 1);
7 }
```

↙
1 · 2 · 3 · 4 · ... (n-1) · n
→

```
1 public static int factorialTail(int n) {
2     return factorialTailHelper(n, 1);
3 }
4
5 private static int factorialTailHelper(int n, int result) {
6     if (n == 1) return result;
7     return factorialTailHelper(n - 1, n * result);
8 }
```

f

f

In this case, we still have $\mathcal{O}(n)$ multiplications (*), but a compiler might perform *tail call optimization* (TCO), eliminating the need for a new stack frame (just not in **Java**).

See you on Thursday!

- Lab 4 (Bucket Sort) due tonight at 11:59pm.
- Start Homework 4!
 1. Implement Radix Sort.
 2. Derive number of `=` for `DIYList.add` method with a different growth technique.
- Reminder that Noah ([go/noah](#)) and Smith ([go/smith](#)) have office hours throughout the week and the 201 Course Assistants have drop-in hours in the late afternoons/evenings ([go/cshelp](#)).
- Submit exit ticket 5T today.