



Middlebury

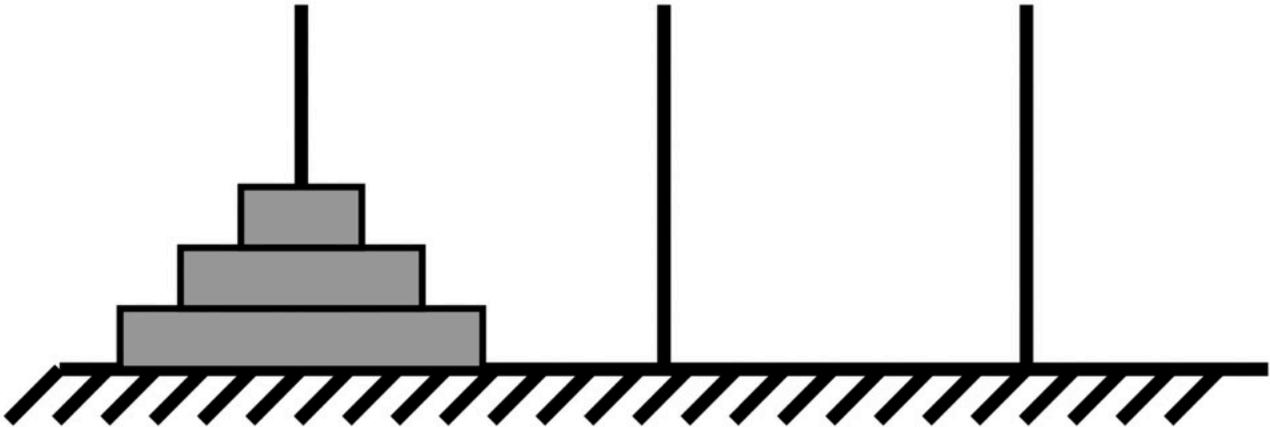
CSCI 200: Math Foundations of Computing

Spring 2026

Lecture 4M: Recurrence Relations I

Goals for today:

- Develop recurrence relations to analyze the performance of recursive algorithms.
- Solve recurrence relations using the *expand and pray* method.
- Solve homogeneous linear recurrence relations.



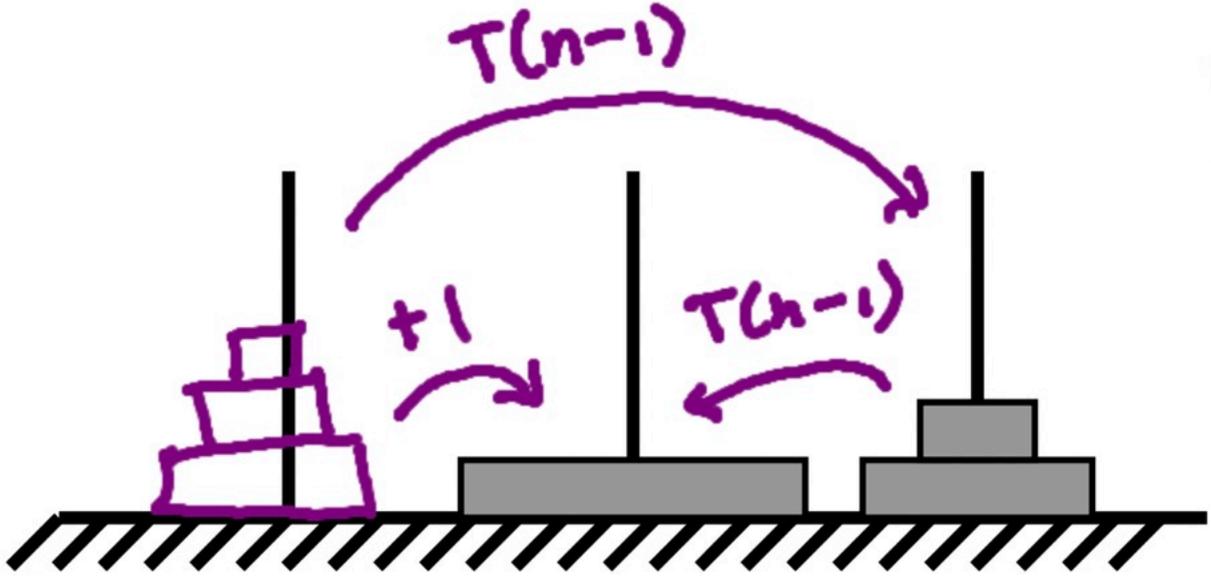
	# moves	
$n = 1$	1	
$n = 2$	3	
$n = 3$	7	
$n = 4$	17	15
	?	

Rules:

1. Move one disk at a time.
2. Every move displaces a disk from the top of one stack to the top of another (or empty rod).
3. No larger disk can ever be placed on top of a smaller one.



Counting the number of moves in the Tower of Hanoi problem.



let $T(n)$ be
moves to
displace a
stack of n
disks

$$T(n) = 2T(n-1) + 1$$

recurrence
relation

base
case

$$T(1) = 1$$

How to get a closed-form expression? Let's try expanding the recurrence relation, and hope we see a pattern.

$$T(n) = 2T(n-1) + 1 \quad \rightarrow \text{plug in } T(n-1) = 1 + 2T(n-2)$$

$$= 1 + 2T(n-1)$$

$$= 1 + 2(1 + 2T(n-2))$$

$$= 1 + 2 + 2^2 T(n-2)$$

$$= 1 + 2 + 2^2 [1 + 2T(n-3)]$$

$$= 1 + 2 + 2^2 + 2^3 T(n-3)$$

$$= 1 + 2 + 2^2 + \dots + 2^i T(n-i)$$

$$= \sum_{i=0}^{n-1} 2^i = \frac{1 - 2^{n-1+1}}{1-2} = \frac{1 - 2^n}{1-2} = 2^n - 1$$

$i = n-1$

$T(n-2) = 1 + 2T(n-3)$

eventually going to reach $T(1) = 1$

when is $n-i = 1$?

last term $i = n-1$

geometric series



Types of recurrence relations we will see.

```
1 def fib(n):  
2     if n <= 1:  
3         return n  
4     return fib(n - 1) + fib(n - 2)
```

linear

$$f(n) = f(n-1) + f(n-2)$$

$$f(0) = 0$$

$$f(1) = 1$$

```
1 def binary_search(lst, val, lo, hi):  
2     if lo > hi:  
3         return -1  
4     m = (lo + hi) // 2  
5     if lst[m] == val:  
6         return m  
7     elif lst[m] > val:  
8         return binary_search(lst, val, lo, m - 1)  
9     else:  
10        return binary_search(lst, val, m + 1, hi)
```

divide-and-conquer

Developing recurrence relations from recursive programs.

```
1 def fib(n):  
2     if n <= 1:  
3         return n  
4     return fib(n - 1) + fib(n - 2)
```

$$f(n) = f(n-1) + f(n-2)$$

$$f(0) = 0$$
$$f(1) = 1$$

homogeneous recurrence relation

Solving homogeneous linear recurrence relations:

Determining a closed-form expression for Fibonacci numbers.

$$f(0) = 0$$
$$f(1) = 1$$

$$f(n) = f(n-1) + f(n-2)$$

$$r^n = r^{n-1} + r^{n-2}$$

$$r^n - r^{n-1} - r^{n-2} = 0$$

$$\begin{bmatrix} r^2 & -r & -1 \end{bmatrix} = 0$$

$\neq 0$ $= 0$

$$r = \frac{1 \pm \sqrt{1+4}}{2}$$

$$r_1 = \frac{1 + \sqrt{5}}{2}, \quad r_2 = \frac{1 - \sqrt{5}}{2}$$

guess: always r^n
 $r \in \mathbb{R}$

$$f(n) = r^n$$

$$f(n-1) = r^{n-1}$$

$$f(n-2) = r^{n-2}$$

thm: $f_1(n) = r_1^n$

both solutions $\rightarrow f_2(n) = r_2^n$

then $c_1 f_1(n) + c_2 f_2(n)$
is also a solution.
 $c_1, c_2 \in \mathbb{R}$

$$f(n) = c_1 \left(\frac{1 + \sqrt{5}}{2} \right)^n + c_2 \left(\frac{1 - \sqrt{5}}{2} \right)^n$$



Solving homogeneous linear recurrence relations:
Determining a closed-form expression for Fibonacci numbers.

$$f(n) = c_1 \left(\frac{1+\sqrt{5}}{2} \right)^n + c_2 \left(\frac{1-\sqrt{5}}{2} \right)^n \quad \text{with } f(0)=0 \\ f(1)=1$$

$$f(0) = 0 = c_1 + c_2$$

$$f(1) = 1 = c_1 \left(\frac{1+\sqrt{5}}{2} \right) + c_2 \left(\frac{1-\sqrt{5}}{2} \right)$$

$$c_1 = \frac{1}{\sqrt{5}} \quad c_2 = -\frac{1}{\sqrt{5}}$$

0, 1, 1, 2, 3, 5, 8, ...

$$f(n) = \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^n$$

What if we have this recurrence relation?

$$f(n) = 4f(n-1) - 4f(n-2), \quad \text{with } f(0) = 1, \quad \text{and } f(1) = 0$$

What are the roots of the characteristic equation for this recurrence?

A. $r_1 = 2, r_1 = -2$

B. $r_1 = 2, r_1 = 4$

C. $r_1 = 2, r_2 = 2$

D. $r_1 = 2, r_1 = 1$

Characteristic equation:
 $r^2 + \dots$

$$r^n = 4r^{n-1} - 4r^{n-2}$$

$$r^n - 4r^{n-1} + 4r^{n-2} = 0$$

$$\underbrace{r^{n-2}}_{\neq 0} \left[\underbrace{r^2 - 4r + 4}_{=0} \right] = 0$$

$r_1 = 2$
 $r_2 = 2$

$$(r-2)^2 = 0$$

thm: repeated roots
 if $f(n) = r^n$
 with r is a
 repeated root
 then $f(n) = n r^n$
 is also a solution.

$$f(n) = c_1 2^n + c_2 \cdot n \cdot 2^n$$

$$f(n) = 2^n$$

$$f(0) = 1 = c_1$$

$$f(1) = 0 = 2 + c_2(1)(2) \rightarrow c_2 = -1$$

$$f(n) = 2^n - n 2^n = (1-n)2^n$$

