



Middlebury

CSCI 200: Math Foundations of Computing

Spring 2026

Lecture 1W: Logic

Last time, we introduced proofs.

A mathematical proof is a verification of a **proposition** by a chain of logical deductions from a set of **axioms**.

Theorem: Given a planar right triangle with side lengths a , b and c (meaning there is one angle of 90° , which we will take to be opposite the side with length c), then $c^2 = a^2 + b^2$.

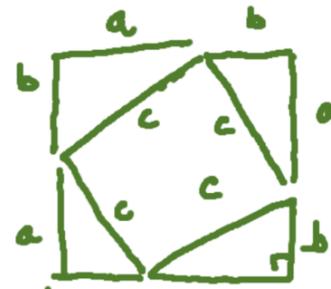
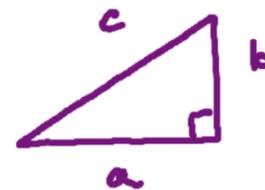
Proof:

$$\Delta_{abc} = \frac{1}{2} ab$$

$$\square_c = c^2$$

$$\square_{a+b} = (a+b)^2 = a^2 + 2ab + b^2$$

$$\square_{a+b} = \square_c + 4 \Delta_{abc} = a^2 + 2ab + b^2$$
$$= c^2 + 4 \left(\frac{1}{2} ab \right) = c^2 + 2ab$$



$$\hookrightarrow a^2 + b^2 = c^2$$

QED



Goals for today:

By the end of today's lecture, you will be able to:

- Create **compound statements** with \wedge , \vee and \neg , and implications (\implies , \iff).
- Prove compound statements using a **truth table**.

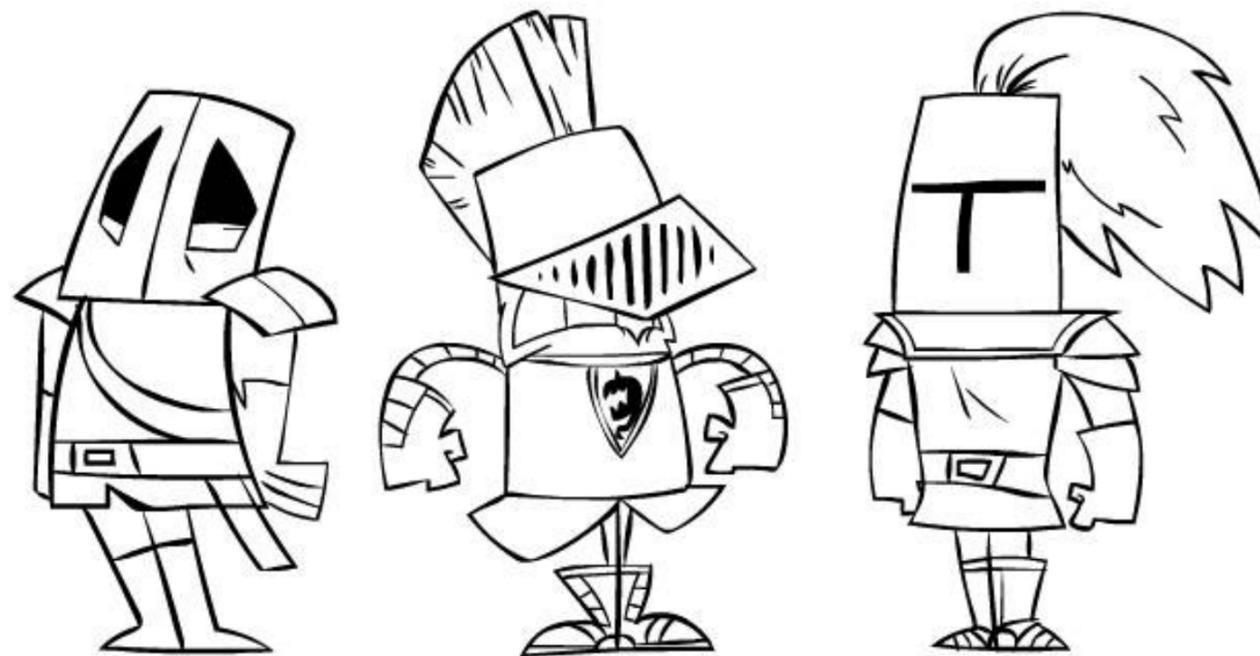


Image credit

But first! A logic puzzle...

You encounter three trolls which are either knights (always tell the truth) or knaves (always lie). They let you pass if you can identify who is a knight and who is a knave. Let's just start with Troll A. Can Troll A be a knight?

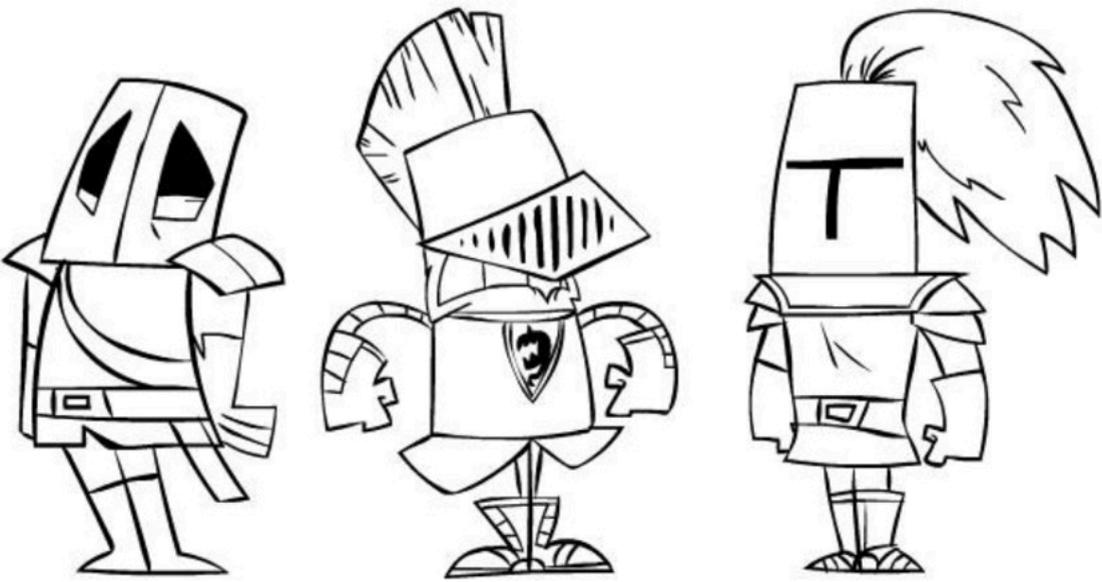
F ?

Suppose A is a knight

- Troll A: Either we are all knaves or two of us are knights.
- Troll B: If I am a knight, then the other two are knaves.
- Troll C: Troll B is lying.

Troll B be a knight? X

Troll B be a knave? X



Troll A is a knave

Logical operators: \wedge , \vee , \neg

Let $p = I \text{ am wearing a blue shirt today.}$

Let $q = \text{Today is Monday.}$

$p \wedge q$: I am wearing a blue shirt AND today is Monday.

$p \vee q$: I am wearing a blue shirt OR today is Monday

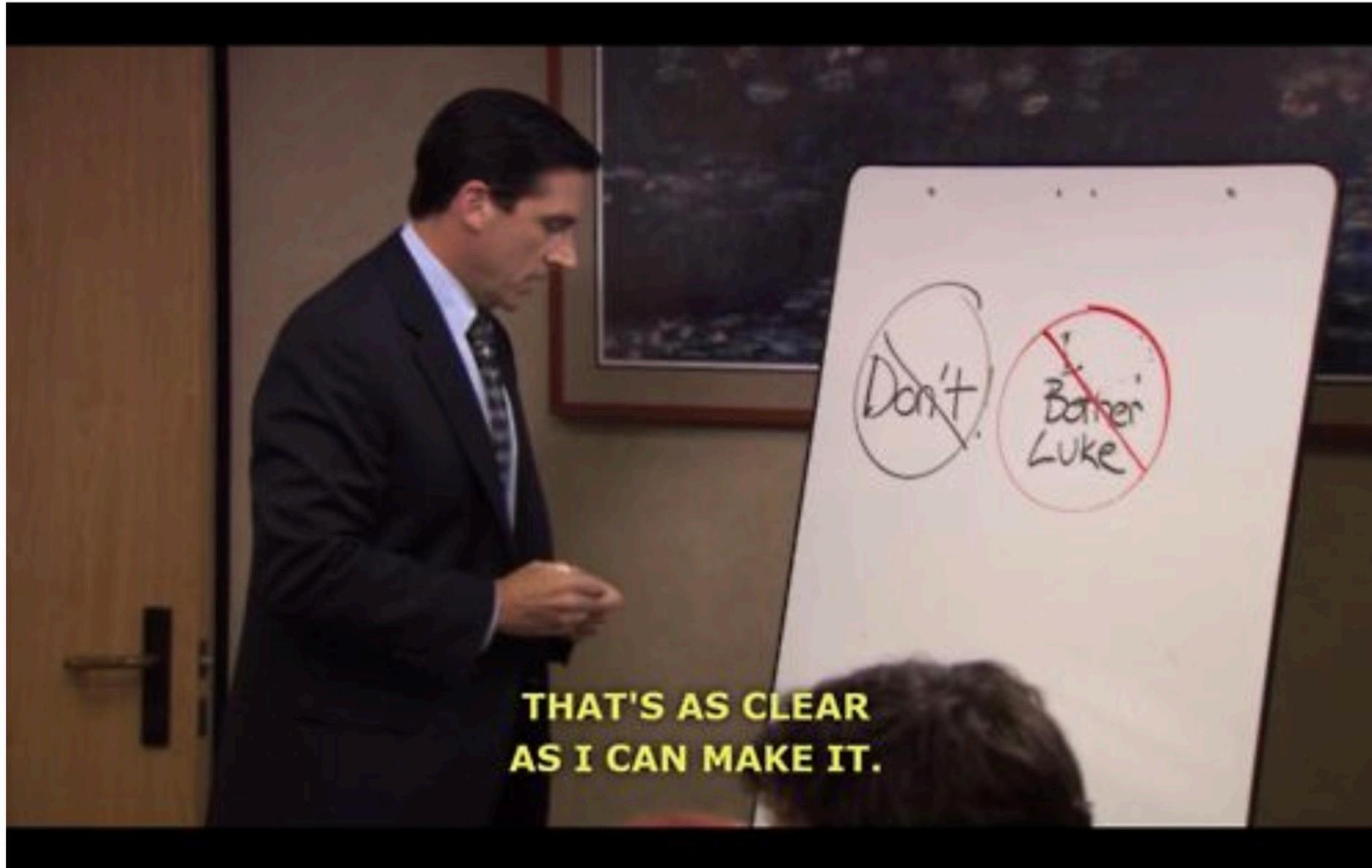
$\neg p$: I am NOT wearing a blue shirt

p	q	$p \wedge q$	$p \vee q$	$\neg p$
T	T	T	T	F
T	F	F	T	F
F	T	F	T	T
F	F	F	F	T



This notation is useful to avoid situations like this...

\wedge AND
 \vee OR
 \neg NOT



~~no, not, u~~



Implications: think of a "promise".

Let $p =$ You get 100% on the final.

Let $q =$ You get an A in the course.

"if p , then q "

$$P \rightarrow Q$$

\rightarrow is the promise valid

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

Is this statement True or False? Vote at *Reactions* page on course website.

If $1 + 1 = 3$ then I am a dinosaur.

- A. True
- B. False



A	B	C	D	E	👏	❤️	👍	👎	😬	🤔	😊	🐧	🐢	🐍
26	8	0	0	0	0	0	2	0	0	0	0	0	9	0



Biconditional: "if-and-only-if"

$$p \iff q = (p \rightarrow q) \wedge (q \rightarrow p)$$

p	q	$p \rightarrow q$	$q \rightarrow p$	$p \iff q$
T	T	T	T	T
T	F	F	T	F
F	T	T	F	F
F	F	T	T	T

Proving with a truth table.

- Identify and label smaller propositions (with variables).
- Write overall proposition symbolically.
- Make a column for each variable and build up to overall proposition.

$$(P \rightarrow Q) \vee (Q \rightarrow R)$$

2^3
rows

Example: Prove the following: *If you eat spinach everyday, then you will win the lottery, or if you win the lottery, you will lose your job.*

P	Q	R	$(P \rightarrow Q) \vee (Q \rightarrow R)$
T	T	T	
T	T	F	
T	F	T	
T	F	F	
F	T	T	
F	T	F	
F	F	T	
F	F	F	

all rows are true ✓

Back to our original puzzle.

$T \rightarrow T$ is T
 $F \rightarrow F$ is T

$p \leftrightarrow t_1$
 $q \leftrightarrow t_2$
 $r \leftrightarrow t_3$ } all must be true

Let p, q, r represent whether Trolls A, B, C are a knight (True) or a knave (False). We already know $\neg p$ is true (Troll A is a knave).

- **Troll A:** Either we are all knaves or two of us are knights.

$$t_1 = (\neg p \wedge \neg q \wedge \neg r) \vee ((p \wedge q) \vee (p \wedge r) \vee (q \wedge r)) \equiv (\neg q \wedge \neg r) \vee (q \wedge r)$$

- **Troll B:** If I am a knight, then the other two are knaves.

$$t_2 = q \implies (\neg p \wedge \neg r) \equiv q \rightarrow \neg r$$

- **Troll C:** Troll B is lying.

$$t_3 = \neg q \quad t_3 = \neg t_2$$

- **Consistency:**

$$c = (p \iff t_1) \wedge (q \iff t_2) \wedge (r \iff t_3)$$

q	r	t_1	t_2	t_3	c
T	T	F	T	F	T ✓
T	F	F	F	F	F
F	T	F	T	F	F
F	F	F	F	F	F

A: knave
 B: knight
 C: knave



de Morgan's laws: a tool to simplify compound propositions.

$$\neg (p \wedge q) \equiv \neg p \vee \neg q$$


$$\neg (p \vee q) \equiv \neg p \wedge \neg q$$


Summary & Reminders

- Problem Set 1 due (on paper) on Friday at 9:45am: Quiz 1 will ask you to re-solve one of these problems.
 - Because it's the first quiz, I'm telling you which problem I will ask: I am going to ask you to solve Problem 0.8. For (b), why is $x = 7$ not a solution?
 - But you must show all your work: 4/5 for each part if only the answer is provided, but no work is shown.
 - Come visit during office hours if you have questions!
General office hours still TBD, but for now:
 - Tomorrow (Thursday): 1pm - 3pm

