



Middlebury

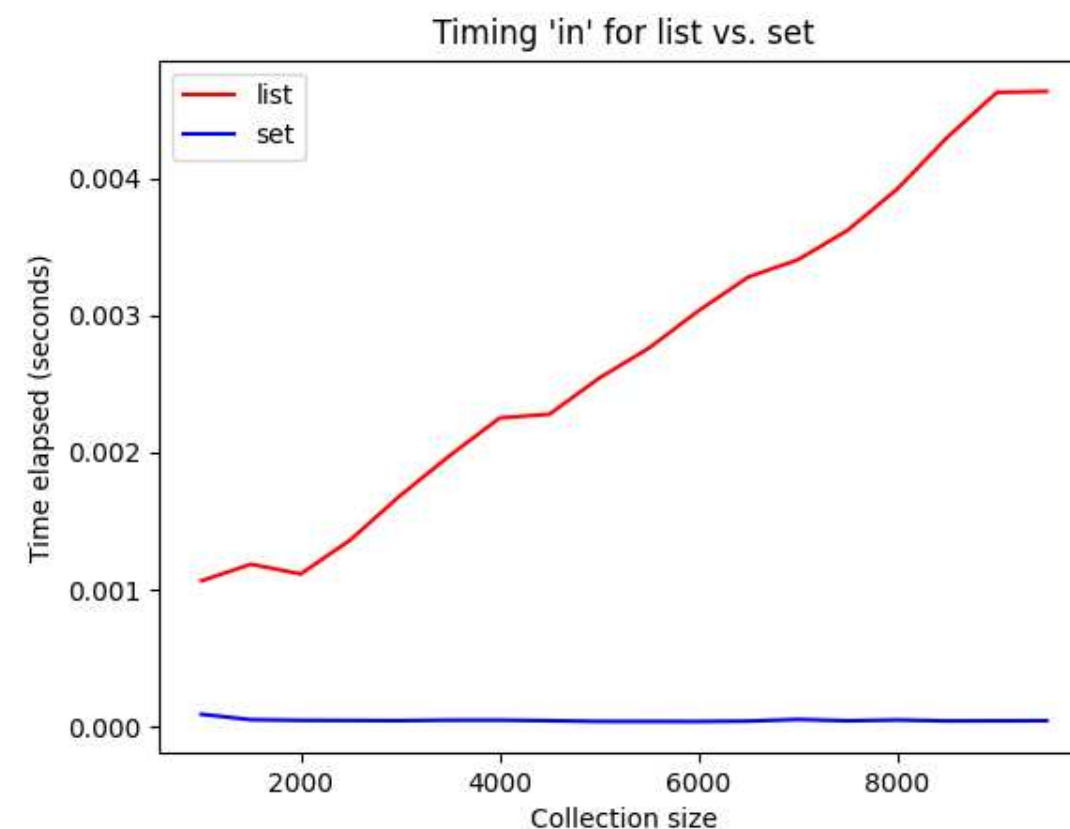
CSCI 146: Intensive Introduction to Computing

Fall 2025

Lecture 16: Complexity analysis and numerical representation

Goals for today

- Informally define **asymptotic complexity** with a focus on **runtime**.
- Describe the purpose and calculation of **Big-O notation**.
- Compare (e.g. rank) constant, logarithmic, linear, quadratic, and cubic complexities.
- Predict the runtime of an algorithm based on its Big-O complexity.
- Explain how integers and other types are represented in the computer.
- Perform binary addition of unsigned integers.
- Be aware of twos-complement representation.
- Be aware of the floating point representation and some of its limitations.



Analyzing the **in** operator.

Possible implementation:

```
def list_in(lst, x):  
    """
```

Should be the same as:

```
>>> x in list
```

```
"""
```

```
for value in lst:
```

```
    if value == x:
```

```
        return True
```

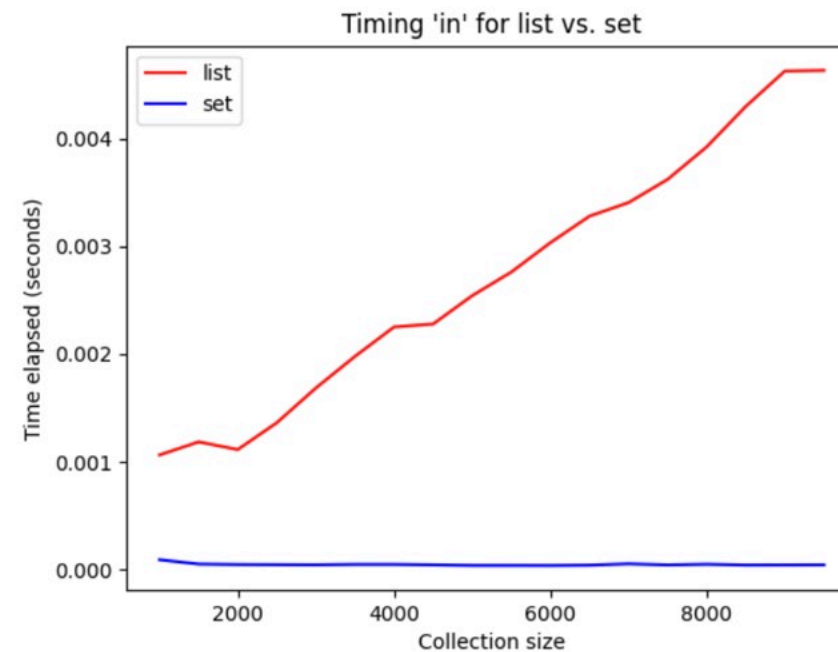
```
return False
```

← "contains"

how many == performed?

$n = \text{len}(\text{lst}) \quad \approx n ==$

$\# \text{ ops} = a n + b$
 ↑ ↑
 constants



Big-O notation allows us to describe the *asymptotic complexity* (runtime, memory) of an algorithm (not a specific implementation).

For the **in** operator on **lists**: (assume a **list** of **len** n)

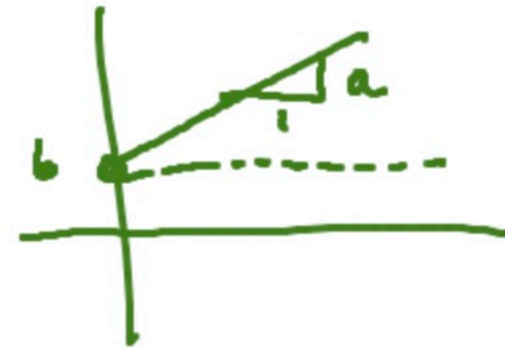
→ $\text{runtime}(n) = an + b$, where a, b are constants.

Main idea: how does the runtime (or memory) grow as the size of the inputs grows? $n \rightarrow \infty$

- Big-O places an *upper-bound* on the growth of the function.
- If the function is a sum of several terms, only the fastest growing term is kept.
- Constant terms (i.e. those independent of the size of the input n) are omitted.

runtime → $f(n) = an + b$

$O(n)$
↑
big-O



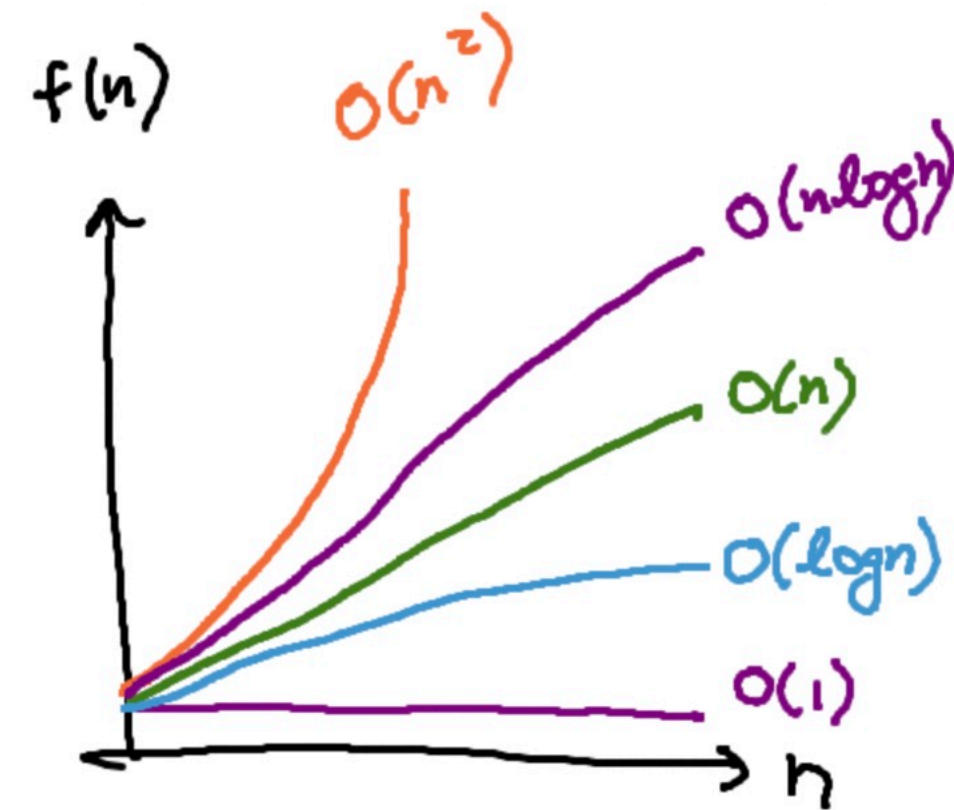
Example: determine a Big-O bound on the following function:

$$f(n) = \underbrace{9n^3}_{\text{dominant}} + \cancel{6n^2} + \cancel{n} + \cancel{5}.$$

$f(n)$ is $O(n^3)$

Common Big-O functions

Complexity	Description
$O(1)$	Constant
$O(\log n)$	Logarithmic
$O(n)$	Linear
$O(\underline{n} \log \underline{n})$	Linearithmic
$O(n^2)$	Quadratic



Estimating runtime using Big-O

Suppose that we had an algorithm with $O(n^2)$ complexity that takes 5 seconds to run on our computer when n is 1000. If the input increased to n of 3000, about how long would it take to run?

$$t = an^2$$
$$t_0 = 5 = a(1000)^2 \quad (1)$$
$$t_1 = a(3000)^2 \quad (2)$$

$$\frac{(2)}{(1)} = \frac{a(3000)^2}{a(1000)^2} = \frac{t_1}{t_0} = 9$$

$$t_1 = 9 \times 5 = \boxed{45 \text{ sec}}$$

What is the complexity of these two implementations of a standard deviation function?

```
def mean(data):  
    """  
    Return mean of iterable data  
    """  
    return sum(data) / len(data)  
  
def stddev(data):  
    """  
    Return standard deviation for iterable data  
    """  
    result = 0.0  
    average = mean(data)  
    for elem in data:  
        result += (elem - average) ** 2  
    return math.sqrt(result / (len(data) - 1))  
  
def stddev2(data):  
    """  
    Return standard deviation for iterable data  
    """  
    result = 0.0  
    for elem in data:  
        result += (elem - mean(data)) ** 2  
    return math.sqrt(result / (len(data) - 1))
```

$n-1$ additions for sum $\rightarrow O(n)$

$O(n) \quad a_1n + b_1$

$a_2n + b_2$

$$\begin{aligned} t &= a_1n + b_1 + a_2n + b_2 \\ &= (a_1 + a_2)n + (b_1 + b_2) \\ &O(n) \end{aligned}$$

$a_2(a_1n + b_1)n + b_2$

$O(n^2)$

Question 1: What is the big-O asymptotic bound of the following function?

$$f(n) = \cancel{3}n^2 + \cancel{4}n + \cancel{2}$$

$O(n^2)$

- A. $O(n)$
- B. $O(n^2)$
- C. $O(n^2 + n)$
- D. $O(3n^2 + 4n + 2)$

A B C D E 🖐️ ❤️ 👍 👎 😬 🤔 😊 🐧 🐢 🐍

0 40 0 0 0 0 4 0 0 0 0 0 0 2 4 1



Question 2: Which of the following functions has the same big-O asymptotic bound as the following function

$$f(n) = 2n^3 + 5n + 2$$

- A. $2n^2 + 5n + 2$
- B. $n^3 + 2n^2$
- C. $n(2n^3 + 5n + 2)$
- D. 2

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0 43 0 0 0 1 0 11 3 0 29 1 55 34 66



Question 3: What is the time complexity of the Python **max** function (example below):

```
>>> numbers = list(range(100))  
>>> the_max = max(numbers)
```

$[0, 1, 2, \dots, 99]$

- A. $O(1)$
- B. $O(100)$
- C. $O(n)$
- D. $O(n \log n)$
- E. $O(n^2)$

```
m = numbers[0]  
for n in numbers:  
    if n > m:  
        m = n  
return m
```

Question 4: The code below implements matrix multiplication between two $n \times n$ matrices (stored as **lists** of **lists**). What is the time complexity of this algorithm?

$n \times n \times n$

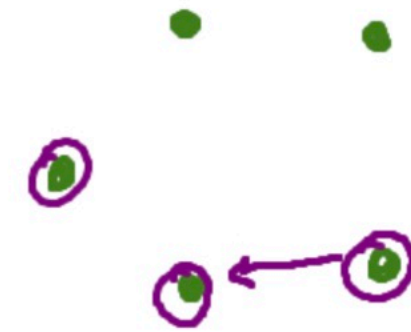
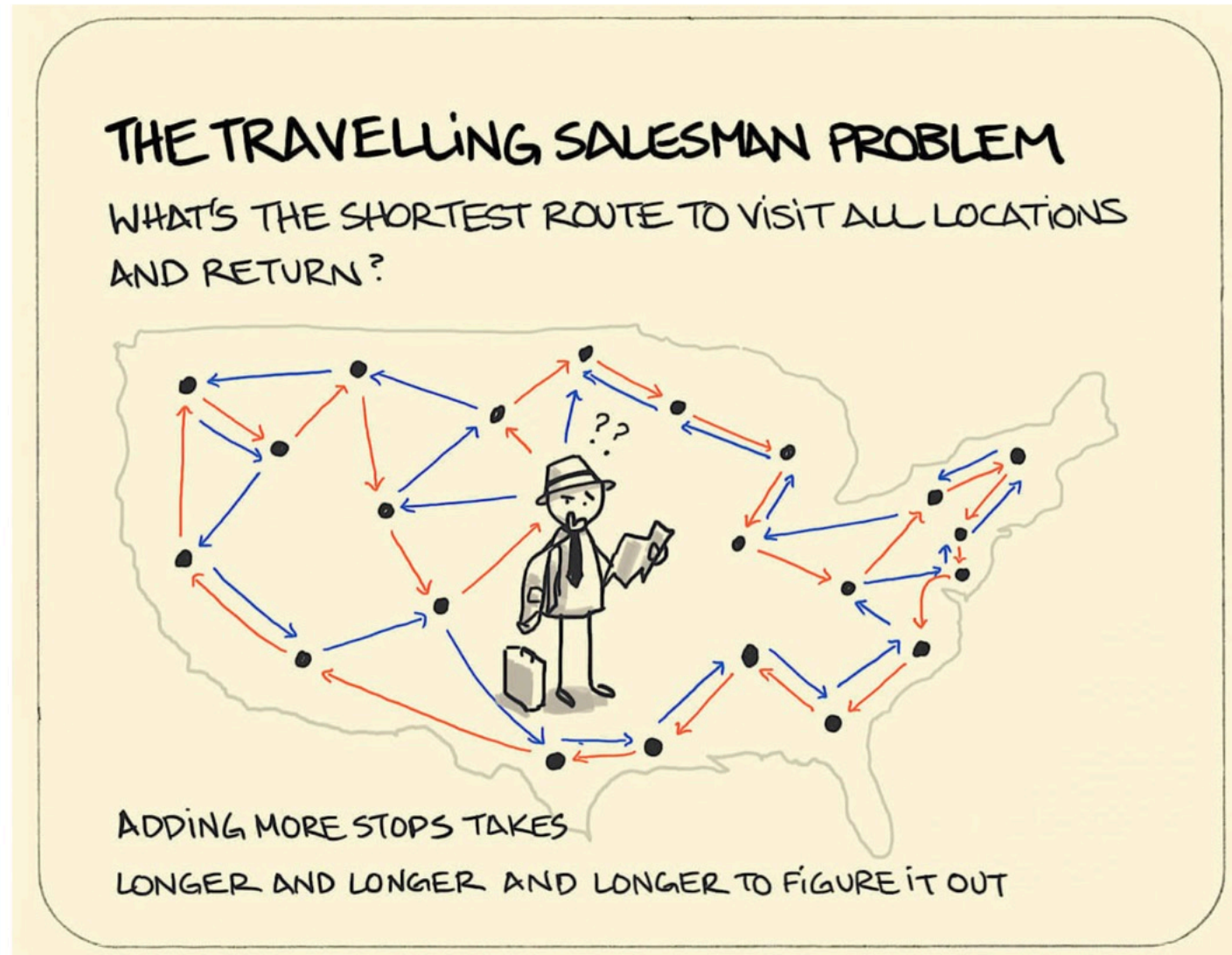
```
for i in range(n):  
    for j in range(n):  
        val = 0.0  
        for k in range(n):  
            val += x[i][k] * y[k][j]  
        res[i][j] = val
```

how many times does this line execute? ❤️

- A. $O(1)$
- B. $O(n)$
- C. $O(n \log n)$
- D. $O(n^2)$
- E. $O(n^3)$

The traveling salesperson problem doesn't have a polynomial time algorithm.

Polynomial time problems: we can find a solution that runs in $O(n^k)$.



$n \times (n-1) \times (n-2) \dots 3 \times 2 \times 1$
↳ all possible paths
 $O(n!)$

source: <https://github.com/ramoneas/travelling-salesman-problem-solver>



The halting problem: can we write a program to determine whether an arbitrary program (with given inputs) *halts*?

Alan Turing proved that such a program does not exist.



```
def will_this_halt():  
    # suppose `halts` solves the halting problem for a given function object  
    if halts(will_this_halt):  
        while True:  
            print("hello from an infinite loop")
```

When analyzing complexity, we can count mathematical, relational and logical operators. How does the computer perform these operations?

```
>>> x = 17
```

Computers use a binary (base-2) representation for numbers, i.e. a sequence of *bits* (binary digits) which are either 0 or 1.

decimal: $17 = 1 \times 10^1 + 7 \times 10^0$

binary: instead of powers of 10, use powers of 2

0	0	1	0	0	0	1	$\rightarrow 1 \times 2^4 + 0 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$ $= 16 + 1 = 17$
---	---	---	---	---	---	---	
2^5	2^4	2^3	2^2	2^1	2^0		
32	16	8	4	2	1		

Converting between binary and decimal.

Converting **1000011** to decimal:

$$\begin{array}{ccccccc} 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 2^6 & 2^5 & 2^4 & 2^3 & 2^2 & 2^1 & 2^0 \end{array}$$

$$64 + 2 + 1 = \boxed{67}$$

Converting **437** to binary:

$$\begin{array}{r} 437 \\ - 256 \\ \hline 181 \end{array} \quad \begin{array}{r} 181 \\ - 128 \\ \hline 53 \end{array}$$

$$\begin{array}{r} 53 \\ - 32 \\ \hline 21 \end{array} \quad \begin{array}{r} 21 \\ - 16 \\ \hline 5 \end{array}$$

$$\begin{array}{ccccccccc} 1 & 1 & 0 & 1 & 1 & 0 & 1 & 0 & 1 \\ \hline 2^8 & 2^7 & 2^6 & 2^5 & 2^4 & 2^3 & 2^2 & 2^1 & 2^0 \\ 256 & 128 & 64 & 32 & 16 & 8 & 4 & 2 & 1 \end{array}$$

$\Rightarrow \text{bin}(437)$

Adding and subtracting with binary arithmetic.

Example: computing $23 + 5$.

$$\begin{array}{r} 1'0'1'1'1 \\ + 00101 \\ \hline 11100 \\ 16 \ 8 \ 4 \rightarrow 28 \end{array}$$

Example: computing $17 - 10$.

$$\begin{array}{r} 0110101 \\ - 1010 \\ \hline 0111 \end{array}$$

$$10 - 01$$

$$\begin{array}{r} 011101 \\ \times 10101 \\ \hline 38 \\ 187 \\ + 38 \\ \hline 125 \end{array}$$

Question 5: What is the minimum number of bits (binary digits) to represent the decimal number 32?

A. 3

B. 4

C. 5

D. 6

E. 7

Question 6: How many distinct numbers can be represented with **5** bits (binary digits) ?

A. 15

B. 16

C. 31

D. 32

E. 63

$$\begin{array}{ccc} \underline{0001} & \underline{0001} & \underline{0001} \\ 2 \times 2 \times 2 \dots & & 2^5 \end{array}$$

Question 7: What is the result of adding **01101** and **00100**?

A. **01001**

B. **01101**

C. **10000**

D. **10001**

E. **10101**

$$\begin{array}{r} 1111 \\ + 001000 \\ \hline 10001 \end{array}$$

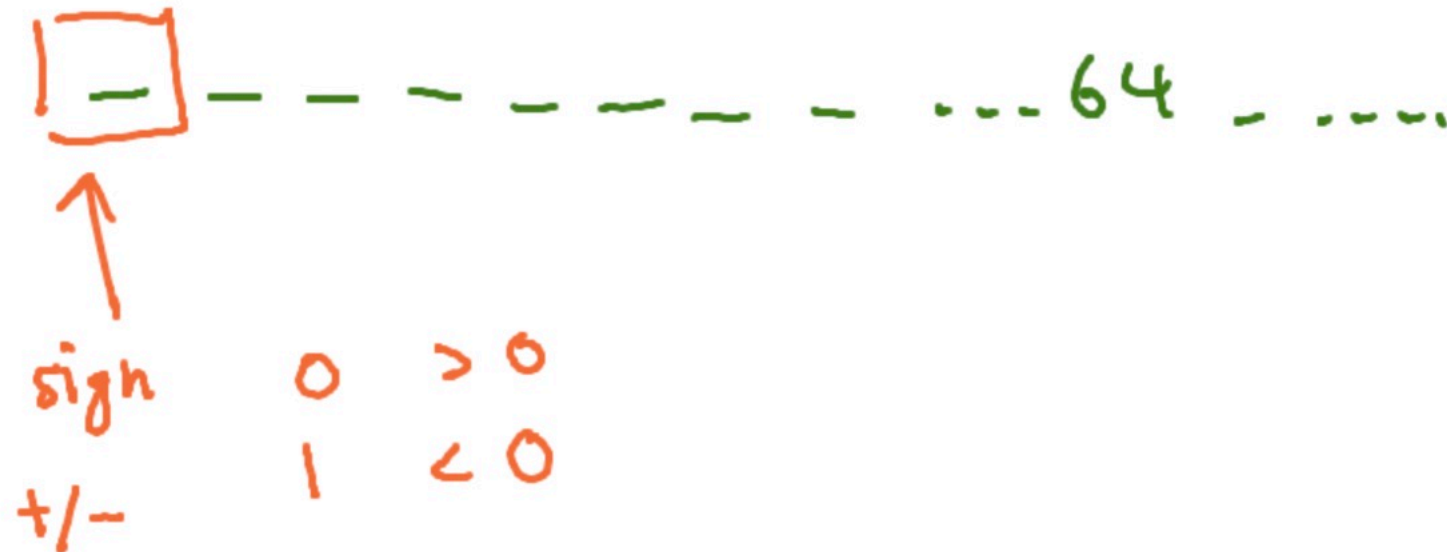
Python uses **sys.maxsize** as the maximum size of a container (and hence maximum index).

On most modern computers, this max size is represented by an integer with 64 bits. So what's **sys.maxsize**?

```
>>> import sys
>>> sys.maxsize
9223372036854775807
```

$\leftarrow 2^{63} - 1$

negative numbers.



Summary and Reminders

- Big-O is concerned with finding an upper bound for our functions:
 1. Find fastest growing terms.
 2. Ignore constants.
- More Big-O analysis applied to searching and sorting algorithms on Wednesday.
- Use the built-in `bin` function to convert a decimal number to binary and check your work.
- Be careful with floating-point expressions: don't use equality checks, instead use tolerances.
- Programming Assignment 7 initial due date on Thursday.
- Quiz 8 this Friday includes retakes from Quizzes 4 - 7 + new Quiz 8 topics + Midterm 1 retakes of 2-3 questions (TBD).
- Use "Regrade Requests" form on the website. See Gradescope comments by clicking on **Code**.